

Recap:

$$1 \xrightarrow{\quad} A \xrightarrow{f_1} E \xrightarrow{f_2} Q \xrightarrow{f_3} 1$$

$$\text{im } f_1 = \ker f_2 : f_1 \text{ inj.}$$

$$\text{im } f_1 = \ker f_2$$

$$\text{im } f_2 = \ker f_3 = Q, f_2 \text{ surj.}$$

1st isomorphism: $E \xrightarrow{f_2} Q \quad \ker f_2 \cong \text{im } f_1$

$$E/\text{im } f_1 \cong \text{im } f_2 = Q$$

2. group actions:

effective: $\forall g \neq 1$. Some x moved.

{

ineffective: $\exists g \neq 1$. All x fixed $\quad \nexists (g, x) = x$

transitive $\quad \forall x, y \quad \exists g, s.t. \quad y = g \cdot x \quad \underline{\text{orb}} = 1$

free: $\forall g \neq 1$ moves all x

3. $\text{Stab}_G(x) = \{g \in G : g \cdot x = x\} \subset G \quad \underline{\text{subgroup}}$

G^x " or isotropy group "

$x^g, \text{Fix}_x(g) = \{x \in X : g \cdot x = x\} \subset X \quad \underline{\text{subset}}$

4. stabilizer-orbit theorem.

$$G/G^x \longleftrightarrow \mathcal{O}_{G^x}$$

$$g \cdot G^x \longleftrightarrow g \cdot x$$

$$y = g_1 \cdot x = g_2 \cdot x \iff g_2^{-1} g_1 \cdot x = x \iff g_2^{-1} g_1 = h \in G^x \iff g_1 = g_2 \cdot h$$
$$\iff g_1 G^x = g_2 h \cdot G^x = g_2 G^x$$

$$\varphi: D_{G^\times} \rightarrow G/G^\times$$

$$g \cdot \infty \mapsto g G^\times$$

$SO(3)$ acts on S^2

$$\text{stab}_{SO(3)}(\hat{\vec{z}}) \cong SO(2)$$

$$"x" = \hat{\vec{z}}$$

$$G^\times = SO(2)_{\hat{\vec{z}}}$$

$$G/G^\times = \underline{SO(3)/SO(2)_{\hat{\vec{z}}}} \cong \text{Dfb}_{SO(3)}(\hat{\vec{z}}) \cup \underline{S^2}$$

$$\pi: SO(3) \rightarrow S^2$$

$$R \mapsto R \cdot \hat{\vec{n}} = \hat{k} \in S^2$$

$$R \cdot \hat{\vec{n}} = R_2 \hat{\vec{n}} = \hat{k} \quad R = R_2 \cdot R_0 \quad R_0 \in \text{stab}(\hat{\vec{n}}) \cong SO(2)$$

"homogeneous space."

7.1 - defn and S-O theorem

7.2 practise with terminologies

3. $SU(2)$ acts on a fiber space \mathbb{C}^2

We know $\vec{x} \in SU(2)$.

$$\vec{x} = \begin{pmatrix} \alpha & \beta \\ -\bar{\beta} & \bar{\alpha} \end{pmatrix} \quad |\alpha|^2 + |\beta|^2 = 1, \quad \alpha, \beta \in \mathbb{C}$$

$$\begin{aligned} \alpha &= x_1 + ix_2 \\ \beta &= x_3 + ix_4 \end{aligned} \quad \rightarrow \quad \sum x_i^2 = 1 \quad \cong S^3$$

Stabilizer-orbit theorem,

$$|\psi\rangle = z_1|0\rangle + z_2|1\rangle$$

$$\vec{z} = \begin{pmatrix} z_1 \\ z_2 \end{pmatrix} \in \mathbb{C}^2 \quad |z_1|^2 + |z_2|^2 = 1$$

$SU(2)$ acts on S^3 transitively

$$\vec{x}(\alpha, \beta, \gamma) = e^{-i\frac{\sigma_x}{2}\gamma} e^{-i\frac{\sigma_y}{2}\beta} e^{-i\frac{\sigma_z}{2}\alpha}$$

$$\hat{z} = |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad \text{stabilizer}$$

$$\begin{pmatrix} \mu & \nu \\ -\bar{\nu} & \bar{\mu} \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} \mu \\ -\bar{\nu} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$\mu = 1, \nu = 0$$

$$stab_{SU(2)}(\hat{z}) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$Dfb_{SU(2)}(\hat{z}) \cong SU(2)/\{ \pm \} = SU(2) \quad SD \text{ theorem}$$

7-3 Centralizer subgroups and counting conj. cls. (69)

Moore

① G acts on G by conjugation

$$O_G(h) = \{g h g^{-1} : g \in G\} =: C(h) \quad \text{conjugacy class}$$

$$\text{Stab}_G(h) = G^h = \{g \in G : g h g^{-1} = h\} =: C_G(h)$$

$(gh = hg)$ centralizer subgroup.

\Rightarrow extend to subset H

$$C_G(H) = \{g \in G : ghg^{-1} = h, \forall h \in H\}$$

$$C_G(H) = Z(G)$$

$$|C(h)| = [G : G^h]$$

② G acts on $X = \{H \subset G : H \neq \emptyset\}$

$$O_G(H) = \{g H g^{-1} : g \in G\}$$

$$G^H = \{g \in G : g H g^{-1} = H\} =: N_G(H)$$

Normalizer

subgroup.

a. $N_G(H)$ is a subgroup.

$$\textcircled{1} e \in N_G(H)$$

$$\textcircled{2} g_1, g_2 \in N_G(H)$$

$$(g_1 g_2^{-1}) H (g_1 g_2^{-1})^{-1} = g_1 (g_2^{-1} H g_2) g_1^{-1}$$

$$= g_1 H g_1^{-1} = H$$

$$\Rightarrow g_1 g_2^{-1} \in N_G(H)$$

~~7.4 Centralizer subgroups & counting conjugacy classes~~

$$|C(h)| = [G : C_G(h)]$$

For a finite group

$$\sum |C(g)| = \frac{|G|}{|C_G(g)|} \quad (\text{stabilizer orbit})$$

$$|G| = \sum_{\substack{\text{distinct} \\ \text{conj. class } \{C(g)\}}} |C(g)| \quad (\text{orbits partition group})$$

$$\Rightarrow |G| = \sum_{\{C(g)\}} \frac{|G|}{|C_G(g)|} \quad \text{"class equation"}$$

Now consider the center

$$Z(G) = \{g \in G : hg = gh \quad \forall h \in G\}$$

$$\forall g \in Z(G), \quad C(g) = \{hgh^{-1} : h \in G\} = \{g\}$$

$$|G| = \sum_{g \in Z(G)} |C(g)| + \sum_{\text{others}} |C(g)|$$

common form

$$= |Z(G)| + \sum_{g \in Z(G)} \frac{|G|}{|C_G(g)|}$$

of class
equation

Theorem. If $|G| = p^n$, p prime, then
center is nontrivial. i.e. $Z(G) \neq \{e\}$

Proof: ① If $C_G(g) = G$. $\exists g \neq e$ trivial.

$\text{Stab}_G(g)$

② Lagrange theorem $\Rightarrow |C_G(g)| = p^{n-u}$ since $n > u$

$$p \mid \sum \frac{|G|}{|C_G(g)|} \Rightarrow p \mid |\Sigma(G)| \text{ i.e. } |\Sigma(G)| \neq 1$$

$= p^n$ ($n > 0$)

Examples $|G| = 8 = 2^3$

Abelian: $\mathbb{Z}_8 \quad Z(\mathbb{Z}_8) = \mathbb{Z}_8$

Non-abelian: $\mathbb{Q} \quad Z(\mathbb{Q}) = \mathbb{Z}_2$

Theorem (Cauchy)

$p \mid |G|$, p prime $\Rightarrow \exists g \in G$ of order p
 $(g^{p-1} = e)$

[HW] Lemma, G abelian, $p \mid |G|$, p prime
 $\Rightarrow \exists g \in G$ of order p .

By induction:

Lemma: G abelian

$p \mid |G|$, p prime $\Rightarrow \exists g \in G$. of order p

Proof. $|G| = p^m$.

the Lemma holds for $m=1$. since if $|G|=p$.

G is cyclic. as a result of Lagrange theorem

then any element $g \in G$ has order p ($g^p = 1$)

Now suppose for a general $m > 1$. $\exists h \in G$. s.t. h has order t ,

i.e. $h^t = 1$

① if $p \mid t$. then $h^{t/p}$ is of order p .

② else $\langle h \rangle$ is a normal subgroup. ($\because G$ is abelian)

$G/\langle h \rangle$ is an abelian group of order

$$|G|/t = p^m/t \quad (\because |\langle h \rangle| = t)$$

then m/t is an integer smaller than m .

by induction. $G/\langle h \rangle$ has an element

of order p

homomorphism $\phi: G \rightarrow G/\langle h \rangle$ a surjection.

$$g \mapsto g\langle h \rangle$$

If $g \in G$ has order p . then

$$\varphi(g^p) = (g \langle h \rangle)^p = 1_{\langle g \rangle \langle h \rangle} = \langle h \rangle$$

$$g^p = h^x \in \langle h \rangle$$

$$\text{if } h^x = 1 \Rightarrow g^p = 1$$

$$\text{else } \exists y \text{ s.t. } (h^x)^y = 1 \Rightarrow g^{py} = 1 \Rightarrow (g^p)^y = 1$$

Proof. (by induction)

$$|G| = pm \text{ holds for } m=1 \quad \checkmark$$

—

If $g \notin Z(G)$, then $|C_G(g)| > 1$, then
 $\subset |G|?$

① $P \mid |C_G(g)| \Rightarrow C_G(g)$ has an element of order p .

② $P \nmid |C_G(g)| (\forall g \in G) \quad |G| = [G : C_G(g)] \underline{|C_G(g)|}$

$$\Rightarrow P \mid [G : C_G(g)]$$

$$|G| = |Z(G)| + \sum \frac{|G|}{|C_G(g)|}$$

$$\Rightarrow P \mid |Z(G)|$$

$\Rightarrow g \in Z(G)$ of order p .

7.5. Example applications of the stabilizer concept

1. In solid state physics, we talk about "little group":

$$\{ g \in P \mid gk = k + K \} \quad \begin{matrix} \downarrow \\ K \text{ reciprocal lattice} \end{matrix}$$

irreps of little group at k determine the band degeneracy etc.

2. Stabilizer code in Quantum information

(for details and more general error-correcting code: see "QC and QI" by Nielsen & Chuang)

Chapter 10 (10.5)

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

X. Y. Z. gates / Pauli matrices $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

$$X|0\rangle = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = |1\rangle$$

$$X|1\rangle = |0\rangle$$

"bit-flip"

$$Z|0\rangle = |0\rangle$$

"phase-flip"

$$Z|1\rangle = -|1\rangle$$

Consider the Pauli group $P^n = (P_i)^{\otimes n}$

$$P_i = \{ \pm I, \pm iI, \pm X, \pm iX, \pm Y, \pm iY, \pm Z, \pm iZ \}$$

and its group action on the vector space

spanned by n -qubit states.

$$\begin{cases} G = P^n \\ X = (\mathbb{C}^2)^{\otimes n} \end{cases}$$

Define $V_S = \{ |\psi\rangle : \underbrace{S|\psi\rangle = |\psi\rangle}_{\text{for } S \in S} \}$

where $S \subset P^n$ a subgroup.

V_S is the vector space stabilized by S

S is the stabilizer of space V_S .

For V_S to be nontrivial.

1. $\forall S_1, S_2 \in S \quad S_1 S_2 = S_2 S_1 \quad S$ abelian

$$\begin{aligned} S, S_1 |\psi\rangle &= S_1 S |\psi\rangle = S_1 |\psi\rangle = |\psi\rangle \\ S_1, S_2 \end{aligned}$$

2. $\lambda I \in S \quad \lambda I |\psi\rangle = |\psi\rangle \quad \lambda = 1$

i.e. $-I, \pm iI \notin S$

$$(-I |\psi\rangle = |\psi\rangle \Rightarrow |\psi\rangle \neq 0)$$

$$z_1 z_2 : \underbrace{1000}_{>}, \underbrace{1001}_{>}, \underbrace{1100}_{>}, \underbrace{1111}_{>}$$

$$z_2 z_3 : \underbrace{1000}_{>}, \underbrace{1100}_{>}, \underbrace{1011}_{>}, \underbrace{1111}_{>}$$

$$V_S = \text{span} \{ 1000, 1111 \}$$

Error set: $\langle x_1, x_2, x_3 \rangle$

$$\{x_1, z_1\} = 0$$

If E anticommutes with $s \in S$

$$s|_{\psi} = 1_{\psi}$$

$$s \underbrace{E|_{\psi}}_{= -E s|_{\psi}} = -E s|_{\psi} = -|\psi\rangle \quad E|\psi\rangle \in V_S^{\perp} \quad \text{measurable by projective measurement}$$

\Rightarrow detectable

If E commutes with $s \in S$ ($E \in N(S) \sim S$)

$$N(S) = \{ fEP^n : fs = sf, \forall s \in S \}$$

$$s \underbrace{E|_{\psi}}_{= E s|_{\psi}} = \underbrace{E|_{\psi}}_{= E|\psi\rangle} \quad E|\psi\rangle \in V_S$$

\Rightarrow undetectable

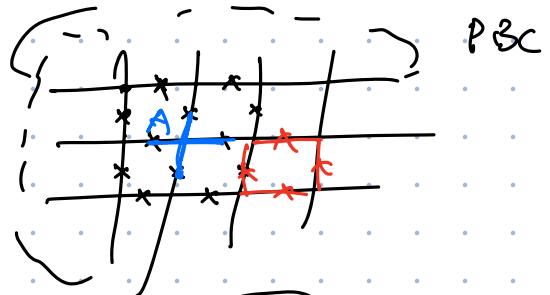
Toric code (Kitaev, Ann. Phys. 2006)

$$A_v = \prod_{j \in \text{star } v} z_j$$

$$B_p = \prod_{j \in \text{plaq.}} x_j$$

$$H = - \sum_v A_v - \sum_p B_p$$

$$[A_v, B_p] = 0$$



$$\frac{A}{B}$$



$S = \langle \{A_v\}, \{B_p\} \rangle$ stabilize the code space V_S

$$N \text{ u.c. } 2^{2N}$$

$$A|\psi\rangle = |\psi\rangle$$

$$A_v^2 = B_p^2 = 1$$

$$B_p|\psi\rangle = |\psi\rangle$$

every A/B cuts the space in half

$$2N \text{ operators, } + \quad \pi A = \pi B = 1$$

(only $N-1$ A/B independent)

$\Rightarrow 2(N-1)$ constraints

$$2^{2N-(2N-2)} = \frac{2^2}{2} = 4$$

ℓ -bit. k -independent generators of S

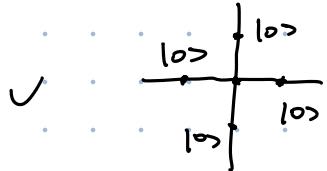
$$\dim(S) = 2^{\ell-k}$$

$$\ell = 2N$$

$$k = 2N-2$$

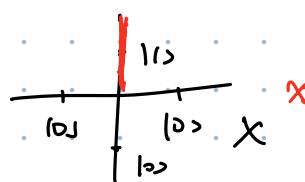
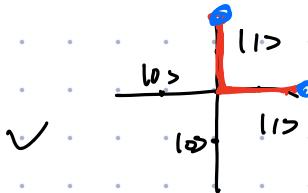
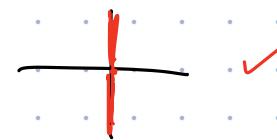
\Rightarrow Toric code encodes two qubits.

A:

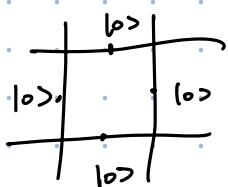


$$Z|0> = |0>$$

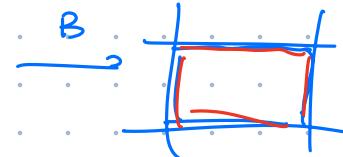
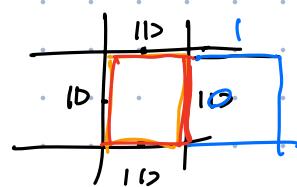
$$Z|1> = -|1>$$



$$B = \pi X$$



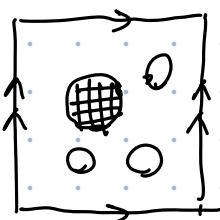
B



equal weight

$\Rightarrow GS = \underline{\text{superposition of all closed loops}}$

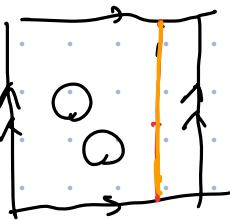
typical config.



$$Z_1 Z_2$$

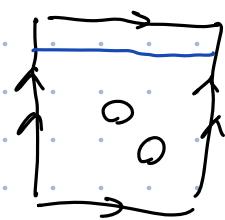
$$(0, 0)$$

$$\downarrow |00>$$



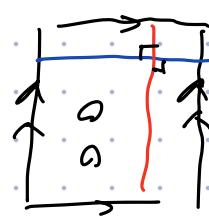
$$(1, 0)$$

$$|10>$$



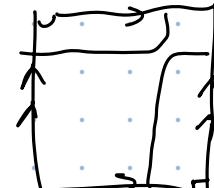
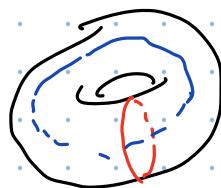
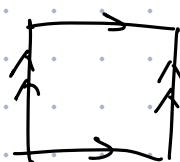
$$(0, 1)$$

$$|01>$$



$$(1, 1)$$

$$|11>$$



local noise/error TZ, πX , suppressed.

Bit operations via stay operators across
the lattice

$$GS = \frac{1}{\sqrt{2}} \left(|B_1\rangle + |B_p\rangle \right) + f |D\rangle + t |D\rangle$$

The diagram illustrates the quantum state GS as a superposition of four components. The first component is $\frac{1}{\sqrt{2}}(|B_1\rangle + |B_p\rangle)$, represented by two adjacent rectangular boxes labeled B_1 and B_p . The second component is $f|D\rangle$, represented by a single rectangular box labeled D . The third component is $t|D\rangle$, represented by a dashed rectangular box labeled D .